

Grid Load Balancing with Schedule Based Backfilling Policy using Gap Search Optimization

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Abstract - Grid computing, one of the most trendy phrase used in IT, is emerging vastly distributed computational paradigm. A computational grid provides a collaborative environment of the hefty number of resources capable to do high computing performance to reach the common goal. Grid computing can be called as super virtual computer, it ensemble large scale geographically distributed heterogeneous resources. Resource allocation is a key element in the grid computing and grid resource may leave at anytime from grid environment. Despite a number of benefits in grid computing, still resource allocation is a challenging task in the grid. This work investigates to maximize the profits by analyzing how the tasks are allocated to grid resources effectively according to quality of service parameter and gratifying user requisition. A method of schedule based backfilling load balancing with gap search algorithm has introduced to answer the above raised question about the resource allocation problem based on grid user requisition. The swift uses genetic algorithms heuristic functions and makes an effective resource allocation process in grid environment. The results of proposed technique of SBBFP-GS algorithm ameliorate the grid resource allocation.

Keywords - Grid Computing, Heterogeneous Resource, Schedule Based Backfilling Load Balancing, Gap Search, Grid.

I. INTRODUCTION

Grids are distributed computational systems that allow access resources owned by organizations. Grid scheduling, that is, the allocation of distributed computational resources to user applications, is one of the most challenging and complex task in Grid computing. Nowadays, several are the applications in which Grids are involved; some practical fields are protein folding, weather modeling, and satellite image processing. In grid architecture, users submit requests for task execution from any one of a number of sites. At each site, besides the local computing system, the system model is composed by three components: an External Scheduler (ES), responsible for determining a particular site where a submitted task can be executed; a Local Scheduler (LS), responsible for determining the order in which tasks are executed at that particular site; a Dataset Scheduler (DS), responsible for determining if and when to replicate data and/or delete local files. Resource site contains, in general, heterogeneous computing resources interconnect by vendor independent networks. In general, on receipt of a task request, the ES interrogates

the LSs to ascertain whether the task can be executed on the available resources and meet the user specified due date. If this is the case, specific site in which executing that task is chosen. Otherwise, the ES attempts to locate LS of a site, controlled by another ES that can meet the task processing requirements, through search mechanisms. If a LS cannot be located within a preset number of search steps, the task request is either rejected or passed to another LS that can minimize the due date failure depending on a task request parameter. When a suitable site is located, the task request is passed from the ES to this site and is managed by the associated LS.

In Following Chapter 2 gives literature reviews of the different method of load balancing, this entire scheme entitled with their authors name and respective title. Chapter 3 describes the process of proposed method with algorithm and corresponding diagrams. Chapter 4 describes the implementation details and information about used data set. The implementation details are also explained out in terms of the algorithm process. Chapter 5 describes the result and analysis of the proposed work in terms of performance metrics. Chapter 6 describes the conclusion of the proposed work and also describes the future work. Chapter 8 specifies the references of research paper, which use the corresponding details in our work.

II. LITERATURE SURVEY

Gnanasekaran et. al[1], Grid Computing emerged as a wide scale distributed system to offer dynamic coordinated resources sharing and high performance computing. Grid coordinates the resources that are not subject to centralized control. It uses standard open, general purpose protocols and its interfaces. Grid delivers non-trivial qualities of service such as the response time, throughput, availability and security.

Caramia et al. [2], Grid scheduling, that is, the allocation of distributed computational resources to user applications, is one of the most challenges and complex task in Grid computing. The problem of allocating resources in Grid scheduling requires the definition of a model that allows local and external schedulers to communicate in order to achieve an efficient management of the resources themselves.



Keerthika et. al[3], Grid Computing provides seamless and scalable access to wide-area distributed resources. Since, computational grid shares, selects and aggregates wide variety of geographically distributed computing resources and presents them as a single resource for solving large scale computing applications, there is a need for a scheduling algorithm which takes into account the various requirements of grid environment.

III. PROPOSED METHODOLOGY

The basic algorithm of proposed work SBBFP-GS is as follows:

[M'] = SBBFP-GS(Q, M)

// Q is the queue for incoming jobs

// M is a map between jobs and resource nodes

// M' is the updated allocation map

Step 1: Initialize the status of all nodes (N).

Step 2: Initial status=.Previous

Step 3: While jobs=N and N>0 do

Step 4: if Current state is ready to change then

Step 5: Current = Get change state (); //Computation stage

Step 6: Threshold = generate threshold (upper bound, lower bound); //Load Balancing

Step 7: Now Gap Search Optimization technique apply on O.

7.1 Initialize gap search parameter θ_0 on the basis of current state and threshold value, initial point $X_0 \subset Q$ and iteration index k=0.

7.2 Generate a collection of candidate points $V_{k+1} \subset Q$ according to a specific generator and associated sampling distribution.

7.3 Update X_{k+1} based on the candidate points V_{k+1} , previous iterates and algorithmic parameters. Also updates algorithm parameters θ_{k+1} .

7.4 If a stopping criterion is met then stop, otherwise increment k and return to step 7.2.

Step 8: Through Gap Search, find the optimized job set in queue Q.

Step 9: Get first job j from Q.

Step 10: while $j \neq null$

Step 11: $N_i \leftarrow$ the number of nodes required by j;

Step 12: $N_{idle} \leftarrow$ the number of idle nodes;

Step 13: if $N_j \le N_{idle}$ then remove j from Q and dispatch

it to any N_j idle nodes; Updates M accordingly;

if j is not at the head of Q then

insert j into Q_{backfill};

else

 $N_{\text{backfill}} \leftarrow$ the number of nodes running jobs arriving later than j;

if
$$N_i \leq (N_{backfill} + N_{idle})$$
 then

Remove j from Q and dispatch it to N_i idle nodes;

Update M;

Step 14: j<- get the next job from Q and goto step 10.

Step 15: Go to step 3 with decrement list of nodes N.

Step 16: M' = M;

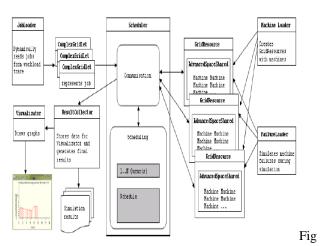
Step 17: Finally obtain the updated jobs and resource node allocation map with required load balancing.

IV. EXPERIMENTAL STEP:

Load Balancing components have been developed which executes in simulated grid environment. This application has been developed using Java and Netbeans. GridSim is event-based modular simulator, composed of independent entities which implement the desired simulation functionality (see Figure 1). It consists of the centralized scheduler, the grid resource(s) with the local job allocation policy, the job loader, the machine and failure loader and additional classes responsible for the simulation setup, the visualization and the generation of simulation output. By now, the Grid users are not directly simulated but the job loader entity can be used as a parent class for the future implementation of the Grid user. Simulator's behavior is driven by the event-based message passing protocol. For each simulated event, such as the job arrival or the job completion, one message defining this



event is created. It contains the identifier of the message recipient, the type of the event, the time when the event will occur and the message data.



1: Main Parts of the GridSim 5.0.2 Simulator

The JobLoader class of GridSim supports several trace formats including the Grid Workloads Format (GWF) of the Grid Workloads Archive and the Standard Workloads Format (SWF) of the Parallel Workloads Archive.

V. RESULT ANALYSIS

The analysis of the existing work (FCFS, EASY-Backfilling) and the proposed work (SBBFP-GS) on different parameters are given in following tables.

Table 1: Average Machine Usage (in Percentage) with Different Load Balancing Strategies

DAYS	FCFS	EASY BACKFILLING	PROPOSED TECHNIQUE (SBBFP-GS)
0-5	13.2	12.8	13.2
5-10	40.6	40.4	42.4
10-15	2.2	1.8	2.2
15-20	45.5	41.6	45
20-25	45.6	44	50
25-30	52	52.4	54.8
30-35	61.4	61	65
35-40	32.4	33	34
40-45	28.6	30	30.8
45-50	57.4	58	62
50-55	60.4	64.2	65.4
55-60	34	34.2	35

We can compare different load balancing technique with our proposed technique. On the basis of load balancing with scheduling method, FCFS (First Come First Serve) and EASY Backfilling used as an existing technique and Schedule Based Backfilling Policy with Gap Search (SBBFP-GS) used as a proposed method. There are four parameters are here for compare results of different approaches.

- (1) Average Machine Usage
- (2) Job Cluster Uses/Day
- (3) Number of Requested and Used CPU
- (4) Number of Waiting and Running Jobs

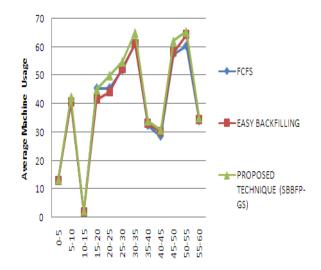


Fig 2: Average Machine Usage (in Percentage) with FCFS, EASY Backfilling and Schedule Based Backfilling with Gap Search Strategies

Table 2: Job Cluster Uses per Day (in Percentage) with Different Load Balancing Strategies

DAYS	FCFS	EASY BACKFILLING	PROPOSED METHOD (SBBFP-GS)
0-5	0.6	0.8	0.6
5-10	40	40	38
10-15	1.8	1.8	1.6
15-20	35.6	31.6	31
20-25	35.2	36.8	34.5
25-30	14.6	15	10.5
30-35	39.8	35.8	35
35-40	12.2	12.6	9.6
40-45	21.4	21.6	20.8
45-50	42.6	44.4	40
50-55	29.8	31	28
55-60	23.8	28.2	23.4

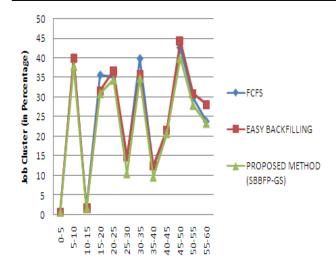


Fig 3: Job Cluster Uses per Day (in Percentage) with FCFS, EASY Backfilling and Schedule Based Backfilling Policy with Gap Search Strategies

As per above graph, proposed technique conservative backfilling with random search has more uses job clusters as compare than FCFS and EASY Backfilling method.

VI. CONCLUSIONS AND FUTURE WORK

The proposed SBBFP-GS Algorithm implements load balancing for scheduling the jobs. Experiments have been done for makespan that serves as a parameter for evaluating the efficiency of the algorithm and finally average resource utilization that serves as the evaluation parameter for proper load balancing. From the results and discussion section it is observed that load balance threshold achieves a better result than the existing FCFS and EASY Backfilling algorithms. The proposed SBBFP-GS algorithm considers the load of each resource at the time of scheduling which are very important in grid environment. This can be extended in future with factors for reducing the communication overhead of the grid system.

We showed the design and implementation of a protocol (SBBFP-GS) for load balancing with scheduling in a Computational Grid. The Grid is partitioned into a number of clusters and each cluster has a coordinator to perform local load balancing decisions and also to communicate with other cluster coordinators across the Grid to provide inter-cluster load transfers, if needed. Our results confirm that the load balancing method is scalable and has low message and time complexities. Our work is ongoing and we are looking into using the proposed model for real-time load balancing where scheduling of a process to a Grid node should be performed to meet its hard or soft deadline.

In future researches nodes can be designed hierarchically and different classes of sites can be considered for nodes (resources) in terms of computational capacity including low, medium and high classes and the efficiency of sites can be discussed based on them. In addition, for evaluating of effectiveness of a node, load of each site can be considered into value function, so that the best site for executing the task can be selected. Future work will focus

- Different scheduling can be applied to optimization techniques
- QoS Constrains such as reliability can be used as performance measure.

Our research in this area is still at an early stage and there are many aspects worthy of further study. First, we have not modeled the impacts of accuracy of job execution time estimation on the effectiveness of our proposed load balancing algorithm. Second, we will utilize migration threshold dynamically based on real-time observation of load behavior of system resources. Finally, we do not take network and hardware failure into account in this study. A failure model may be employed to study this influence. Owing to the dynamic nature of the practical grid environment, designing an ideal load balancing algorithm still remains a challenge.

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